

Formal Semantics in Modern Type Theories

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Structure

- What are MTTs?
- Brief intro to MTTs
- Examples of using MTTs for NL semantics
- Conclusions

Modern Type Theories

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 - ★ Proof theoretic specification

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 - ▶ Alternative foundations for mathematics (Homotopy Type Theory) [Voevodsky(2015)]
 - ▶ Formalization using proof assistants: systems implementing constructive type theories that help in the formalization of mathematics and program verification
 - ★ prime examples: Agda [Agda 2008()], Coq [Coq 2007()]

Modern Type Theories and linguistic semantics

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 - ▶ Many more after that
[Boldini(2000), Cooper(2005), Dapoigny and Barlatier(2009), Bekki(2014), Retoré(2013), Grudzinska and Zawadowski(2014), Chatzikyriakidis and Luo(2012), Chatzikyriakidis and Luo(2017a)] among others
 - ▶ How they are useful and in what ways they are different from STT?

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- In MTTs, no such restriction exists: the universe of entities CAN be many-sorted
 - ▶ Arbitrary number of types can be available giving more structure to the domain of individuals, e.g. *man*, *chair*: *Type* (this is the approach by Ranta, Boldini, Luo and colleagues among others)
 - ★ This is known as the CNs-as-Types approach [Chatzikyriakidis and Luo(2016 (to appear).)]
 - ★ However, this is a choice! Other researchers like Bekki and colleagues working on MTTs, prefer to interpret CNs more standardly, i.e. as predicates [Bekki(2014)]

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 - MS man: $e \rightarrow t$
 - MTTs man: Type

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 - ▶ Talk: *human* \rightarrow *Prop*
 - ▶ the ham: *ham* (with *ham: Type*)
 - ▶ Functional application not possible!

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 - ★ *walk:Animal* \rightarrow *Prop*
 - ★ *the_man:Man* (with *man:Type*)
 - ★ Fine if *man* \leq *human*

Different Systems of Subtyping

- Classic case: Subsumptive subtyping

$$\frac{a:A, A \leq B}{a:B}$$

- a term of type A can be used in a context where a term of type B is required instead just in case $A \leq B$

Different Systems of Subtyping

- Record Type Subsumption: a type of subsumptive subtyping for TTR

$$\left[\begin{array}{l} x : Man \\ y : Donkey \\ e : own(x,y) \end{array} \right]$$

will also be of type

$$\left[\begin{array}{l} x : Man \\ y : Donkey \end{array} \right]$$

and also of type

$$\left[x : Man \right]$$

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 - ★ A is a (proper) subtype of B ($A < B$) if there is a unique implicit coercion c from type A to type B
 - ★ An object a of type A can be used in any context $\mathcal{C}_B[_]$ that expects an object of type B : $\mathcal{C}_B[a]$ is legal (well-typed) and equal to $\mathcal{C}_B[c(a)]$.

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- Metatheoretically more advantageous: canonicity is preserved
 - ▶ Long story!

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- MTTs offer a range of other more advanced typing structures
 - ▶ Dependent Typing
 - ★ A family of types that may depend on some value

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 - ▶ When A is a type and P is a predicate over A , $\Pi x:A.P(x)$ is the dependent function type that stands for the universally quantified proposition $\forall x:A.P(x)$
 - ▶ Π for polymorphic typing: $\Pi A:CN.(A \rightarrow Prop) \rightarrow (A \rightarrow Prop)$

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 - ▶ When A is a type and P is a predicate over A , $\Pi x:A.P(x)$ is the dependent function type that stands for the universally quantified proposition $\forall x:A.P(x)$
 - ▶ Π for polymorphic typing: $\Pi A:CN.(A \rightarrow Prop) \rightarrow (A \rightarrow Prop)$
 - ▶ A is a type and B is an A -indexed family of types, then $\Sigma x:A.B(x)$, is a type, consisting of pairs (a, b) such that a is of type A and b is of type $B(a)$.
 - ▶ Adjectival modification as involving Σ types
[Ranta(1994), Chatzikyriakidis and Luo(2017b)]:
[[*heavy_book*]] = $\Sigma x : book.heavy(x)$

Intro to MTTs-Universes

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- ▶ A universe is a collection of (the names of) types into a type (Martin Löf, 1984).
- ▶ Universes can help semantic representations. For example, one may use the universe CN : *Type* of all common noun interpretations and, for each type A that interprets a common noun, there is a name \overline{A} in CN . For example,

$$\overline{man} : CN \quad \text{and} \quad T_{CN}(\overline{man}) = man.$$

In practice, we do not distinguish a type in CN and its name by omitting the overlines and the operator T_{CN} by simply writing, for instance, $man : CN$.

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 - ▶ A universe over which the coordination operator extends

$$\frac{}{PType : Type} \quad \frac{}{Prop : PType} \quad \frac{A : LType \quad P(x) : PType [x:A]}{\prod x:A.P(x) : PType}$$
$$\frac{}{LType : Type} \quad \frac{}{CN : LType} \quad \frac{A : CN}{A : LType} \quad \frac{A : PType}{A : LType}$$

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 - ▶ Various way of thinking about contexts
 - ★ List of variable declarations, where variables stand for proofs of the corresponding assumptions
 - ★ A sequence of type judgements
 - ★ Formally, a context is an expression of the form:
$$\Gamma = x_1 : A_1, x_2 : A_2(x_1), \dots, x_n : A_n(x_1, \dots, x_{n-1})$$
 - ★ A series of types, and a series of proof objects for these types
 - ★ Any type may depend on any of the previous proof objects

Contexts

- They have been used instead of possible worlds for belief intensionality [Ranta(1994), Chatzikyriakidis and Luo(2013)] and also to formalize discourse structure [Ranta(1994), Boldini(2000), Chatzikyriakidis and Luo(2014)]

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A farmer owns a donkey. He loves it.
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- Consider the following discourse:

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- Following the end of the first sentence, we have:

$x_1 : (\Sigma x : \textit{Farmer})(\Sigma y : \textit{Donkey})(\textit{own}(x, y))$

- The pronouns pick variables already declared using the projection operators (π_1 and π_2)

$x_2 : (\textit{love}(\pi_1(x_1), \pi_1(\pi_2(x_1))))$

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 - ▶ For example, from $\Sigma(\textit{man}, \textit{black})$ it does not follow that something is black or that something is a black human
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 - ▶ Regular type for intersective adjectives
 - ★ *handsome:human* \rightarrow *Prop*
 - ▶ polymorphic for subsectives
 - ★ *skilful:ΠA* : CN. (*A* \rightarrow *Prop*)

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- The first projection π_1 of the Σ is declared as a coercion
 - ▶ $\Sigma(\text{man}, \text{black}) \leq \text{Man}$
 - ★ Thus, from black man we can infer man
 - ▶ Subtyping propagates through the constructors: if $A \leq B$ then for all $P: C \rightarrow \text{Prop}$ (with $A, B \leq C$), $\Sigma(A, P) \leq \Sigma(B, P)$
 - ★ This means that: $\Sigma(\text{man}, \text{black}) \leq \Sigma(\text{human}, \text{black})$

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 - ▶ The modification involves an argument which is the class restriction, so $\Sigma(\textit{surgeon}, \textit{skilful}(\textit{surgeon}))$
 - ★ It does not follow that a skilful surgeon is a skilful human:
 $\Sigma(\textit{surgeon}, \textit{skilful}(\textit{surgeon})) \not\Rightarrow \Sigma(\textit{human}, \textit{skilful}(\textit{surgeon}))$
 - ★ $\Sigma(\textit{human}, \textit{skilful}(\textit{surgeon}))$ is not well-typed,
 $\textit{skilful}(\textit{surgeon}) : \textit{surgeon} \rightarrow \textit{Prop}$ and our π_1 is of type *human*

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- Intensional adjectives: alleged
- A belief context: a sequence of judgments a specific human holds
 - ▶ $\Gamma_p = x_1 : A_1, \dots, x_n : A_n(x_1, \dots, x_{n-1})$
 - ▶ We can similarly define allegation contexts
 - ▶ Let $A_N : \text{CN}$ is the interpretation of a common noun N . Then:
alleged N = $\Sigma a : \text{Human}. \Gamma_a(A_N)$
 - ▶ An alleged N: has been alleged by someone that it is an N
 - ★ Belief and allegation contexts can be kept separate

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 - ▶ Then we define:
 - ★ *healthy* = $\lambda x : \text{Human} . \forall h : \text{Health} . \text{Healthy}(h)(x)$
 - ★ *sick* = $\lambda x : \text{Human} . \neg(\forall h : \text{Health} . \text{Healthy}(h)(x))$
- For gradability and gradable adjectives have a look at [Chatzikyriakidis and Luo(2017a)]

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 - ★ Sentence adverbs: $Prop \rightarrow Prop$
 - ★ VP-adverbs: $\Pi A:CN.(A \rightarrow Prop) \rightarrow (A \rightarrow Prop)$
 - ★ Polymorphic type: Depends on the choice of A
 - ★ Given that we are talking about predicates, depends on the choice of the argument
 - ★ $walk:Animal \rightarrow Prop \Rightarrow ADV(walk):(Animal \rightarrow Prop)$
 - ★ $drive:Human \rightarrow Prop \Rightarrow ADV(drive):(Human \rightarrow Prop)$

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 - ▶ John opened the door quickly \Rightarrow John opened the door
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- Non-veridical adverbs do not have this property
 - ▶ John allegedly opened the door \nRightarrow John opened the door

Adverbial Modification: Veridicality

- Define an auxiliary object first, then define the adverb as its first projection
 - ▶ $VER_{Prop} : \Pi v : Prop. \Sigma p : Prop. p \supset v$
 - ▶ $ADV_{ver-Prop} = \lambda v : Prop. \pi_1(VER_{Prop}(v))$
- An adverb like *fortunately* will be defined as:
- $fortunately = \lambda v : Prop. \pi_1(VER_{Prop}(v))$

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- An adverb like *fortunately* will be defined as:
- $fortunately = \lambda v : Prop. \pi_1(VER_{Prop}(v))$
- Consider the following: Fortunately, John went \implies John went
 - ▶ The second component of $(VER_{Prop}(v))$ is a proof of $fortunately(v) \Rightarrow v$
 - ▶ Taking v to be *John went*, the inference follows

Adverbial Modification: Intensional/domain adverbials

- Use of TT contexts in this case as well
 - ▶ *allegedly* = $\lambda P : Prop. \exists p:human, A_p(P)$
 - ▶ Someone has alleged that P (A_p is an agent's allegation context [Chatzikyriakidis and Luo(2017b)])

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- Introduction of intentional contexts: Contexts including the intentions (rather than the beliefs) of an agent. We can use this idea for adverbs like intentionally:
 - ▶ $intentionally = \lambda x : human. \lambda P : [human] . \lambda P:Prop. I_x(P(x)) \wedge \Gamma(P(x))$

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 - ▶ Someone has alleged that P (A_p is an agent's allegation context [Chatzikyriakidis and Luo(2017b)])
- Introduction of intensional contexts: Contexts including the intentions (rather than the beliefs) of an agent. We can use this idea for adverbs like intentionally:
 - ▶ $intentionally = \lambda x : human. \lambda P : \llbracket human \rrbracket . \lambda P:Prop. I_x(P(x)) \wedge \Gamma(P(x))$
- Domain adverbs, e.g. *botanically*, *mathematically*
 - ▶ $botanically = \lambda P : Prop. \Gamma_B P$
- Intensionality without possible worlds

Co-predication and Dot-types

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(4) John picked up and mastered the book.

- A physical and an informational dimension of *book*
 - ▶ The idea is that book is a complex type with both a physical and an informational aspect

Co-predication and Dot-types

- Predicates requiring different kinds of arguments, are used in coordination and applied to the "same" CN argument

(5) John picked up and mastered the book.

- A physical and an informational dimension of *book*
 - ▶ The idea is that book is a complex type with both a physical and an informational aspect
 - ★ We introduce the dot type constructor, forming complex types from simple types
 - ★ To form a dot type $A \bullet B$, its individual components should not share parts
 - ★ E.g. $\text{PHY} \bullet \text{PHY}$ cannot be a dot-type
 - ★ The dot-type is a subtype of its individual parts

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 - ★ The dot-type is a subtype of its individual parts
 - ▶ A scary slide follows!

The rules for dot-types

Formation Rule

$$\frac{\Gamma \text{ valid } \langle \rangle \vdash A : \text{Type} \quad \langle \rangle \vdash B : \text{Type} \quad C(A) \cap C(B) = \emptyset}{\Gamma \vdash A \bullet B : \text{Type}}$$

Introduction Rule

$$\frac{\Gamma \vdash a : A \quad \Gamma \vdash b : B \quad \Gamma \vdash A \bullet B : \text{Type}}{\Gamma \vdash \langle a, b \rangle : A \bullet B}$$

Elimination Rules

$$\frac{\Gamma \vdash c : A \bullet B}{\Gamma \vdash p_1(c) : A} \quad \frac{\Gamma \vdash c : A \bullet B}{\Gamma \vdash p_2(c) : B}$$

Computation Rules

$$\frac{\Gamma \vdash a : A \quad \Gamma \vdash b : B \quad \Gamma \vdash A \bullet B : \text{Type}}{\Gamma \vdash p_1(\langle a, b \rangle) = a : A} \quad \frac{\Gamma \vdash a : A \quad \Gamma \vdash b : B \quad \Gamma \vdash A \bullet B : \text{Type}}{\Gamma \vdash p_2(\langle a, b \rangle) = b : B}$$

Projections as Coercions

$$\frac{\Gamma \vdash A \bullet B : \text{Type}}{\Gamma \vdash A \bullet B <_{p_1} A : \text{Type}} \quad \frac{\Gamma \vdash A \bullet B : \text{Type}}{\Gamma \vdash A \bullet B <_{p_2} B : \text{Type}}$$

Coercion Propagation

$$\frac{\Gamma \vdash A \bullet B : \text{Type} \quad \Gamma \vdash A' \bullet B' : \text{Type} \quad \Gamma \vdash A <_{c_1} A' : \text{Type} \quad \Gamma \vdash B = B' : \text{Type}}{\Gamma \vdash A \bullet B <_{d_1[c_1]} A' \bullet B' : \text{Type}}$$

where $d_1[c_1](x) = \langle c_1(p_1(x)), p_2(x) \rangle$.

$$\frac{\Gamma \vdash A \bullet B : \text{Type} \quad \Gamma \vdash A' \bullet B' : \text{Type} \quad \Gamma \vdash A = A' : \text{Type} \quad \Gamma \vdash B <_{c_2} B' : \text{Type}}{\Gamma \vdash A \bullet B <_{d_2[c_2]} A' \bullet B' : \text{Type}}$$

where $d_2[c_2](x) = \langle p_1(x), c_2(p_2(x)) \rangle$.

$$\frac{\Gamma \vdash A \bullet B : \text{Type} \quad \Gamma \vdash A' \bullet B' : \text{Type} \quad \Gamma \vdash A <_{c_1} A' : \text{Type} \quad \Gamma \vdash B <_{c_2} B' : \text{Type}}{\Gamma \vdash A \bullet B <_{d[c_1, c_2]} A' \bullet B' : \text{Type}}$$

where $d[c_1, c_2](x) = \langle c_1(p_1(x)), c_2(p_2(x)) \rangle$.

Dot-types

John:Man

picked_up:phy \rightarrow *human* \rightarrow *Prop*

mastered:info \rightarrow *human* \rightarrow *Prop*

the: $\Pi A:CN.A$

book:CN

and: $\Pi A.A \rightarrow A \rightarrow A$

Works because *Book* \leq *Phy* • *Info* \leq *Info, Phy*

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- Individuation criteria: John picked up and mastered three books
 - ▶ Every CN carries its own criteria of identity: CNs as setoids

Criteria of Identity/Individuation: CNs as Setoids

- Individuation is the process by which objects in a particular collection are distinguished from one another
 - ▶ Provides us with means to count and a sameness criterion
- In linguistic semantics, individuation is related to the idea that a CN may have its own identity criterion for individuation [Geach(1962)]
- Mathematically, the association of an equivalence relation (the identity criterion) CNs
 - ▶ In constructive mathematics, a set or a type is indeed a collection of objects together with an equivalence relation that serves as identity criterion of that collection

CNs as Setoids

- CNs are not just types
 - ▶ Types plus an identity criterion for that specific CN

$$(7) \quad (A, =)$$

where A is a type and $=: A \rightarrow A \rightarrow Prop$ is an equivalence relation over A

- ▶ The difference between CNs-as-Types and CNs-as-Setoids

$$(8) \quad [human] = Human : Type \text{ (CNs-as-types view)}$$

$$(9) \quad [human] = (Human, =_h) \text{ (CNs-as-Setoids view)}$$

CNs as Setoids

- Consider the following examples and their semantic interpretations:

(10) Three men talk.

(11) Three humans talk.

(12) $\exists x, y, z : \text{Man}. x \neq_M y \ \& \ y \neq_M z \ \& \ x \neq_M z \ \& \ \text{talk}(x) \ \& \ \text{talk}(y) \ \& \ \text{talk}(z)$

(13) $\exists x, y, z : \text{Human}. x \neq_H y \ \& \ y \neq_H z \ \& \ x \neq_H z \ \& \ \text{talk}(x) \ \& \ \text{talk}(y) \ \& \ \text{talk}(z)$

where $\text{MAN} = (\text{Man}, =_M)$ and $\text{HUMAN} = (\text{Human}, =_H)$ are setoids and the identity criterion for men and that for humans are used to express that x , y and z are distinct from each other.

CNs as Setoids

- Necessary to consider the individuation criteria explicitly by using the identity criteria $=_M$ and $=_H$
- The relationship between the **MAN** and **HUMAN** is one where the first inherits the IC from the second

$$(14) \quad (=M) = (=H)|_{Man}$$

CNs as Setoids: Subsetoids

- $A = (A, =_A)$ is a sub-setoid of $B = (B, =_B)$, notation $A \sqsubseteq B$, iff
 - ▶ $A \leq B$ and $=_A$ is the same as $(=_B)|_A$ (the restriction of $=_B$ over A).
- Some examples:

(15) $MAN \sqsubseteq HUMAN$

(16) $(RTable, =_t) \sqsubseteq (Table, = t)$

where $RTable$ is: $\Sigma x: Table.red(x)$ is the domain of red tables and $=_t$ is the equivalence relation representing the identity criterion for tables

CNs as Setoids: Subsetoids

- In restricted domains like *Man* or *RTable*, the identity criteria coincide with those in *Human* and *Table*
- For these cases, one can ignore the IC, i.e. one can use the simpler CNs-as-Types approach
 - ▶ More sophisticated cases like copredication with quantification however need IC

(17) John picked up and mastered three books.

- Double distinctness

(18) John picked up and mastered three books \Rightarrow John picked up three physical objects and mastered three informational objects

CNs as Setoids: Copredication

- Let us split the example into its conjuncts

(19) *Three(Book, PHY, pick up(j)).*

(20) *Three(Book, INFO, master(j)).*

- Note: the CN *book* in 19 refers to a different collection from that referred to by *book* in 20

(21) $\text{BOOK}_1 = (\text{Book}, =_p)$

(22) $\text{BOOK}_2 = (\text{Book}, =_i)$

CNs as Setoids: Copredication

- How the identity criterion for *books* is determined

- ▶ why do we use $=_p$ in 19 and $=_i$ in 20?

- ★ The verb (and its semantics) that determines the identity criterion of the object CN.

$$IC^{N,V} \Rightarrow \begin{cases} =_p & \text{if Dom}(V) = \text{PHY} \\ =_i & \text{if Dom}(V) = \text{INFO} \\ ??? & \text{if Dom}(V) = \text{PHY} \bullet \text{INFO} \end{cases}$$

- ★ In order to deal with the dot-type case, we have to define setoids for dot-types!

CNs as Setoids: Copredication

- Let $A = (A, =_A)$ and $B = (B, =_B)$ be setoids. Then, the *dot-setoid* $A \bullet B$ is defined as follows:
 - ▶ $A \bullet B = (A \bullet B, =_{A \bullet B})$
where $\langle a_1, b_1 \rangle =_{A \bullet B} \langle a_2, b_2 \rangle$ if, and only if, $(a_1 =_A a_2) \vee (b_1 =_B b_2)$.

CNs as Setoids: Copredication

- The semantics for *three*

Let A be a type and $B = (B, =_B)$ a setoid such that $A \leq B$, and $P : B \rightarrow Prop$ a predicate over B :

- $Three(A, B, P) = \exists x, y, z : A. D[B](x, y, z) \ \& \ P(x) \ \& \ P(y) \ \& \ P(z).$
where $D[B](x, y, z) = x \neq_B y \ \& \ y \neq_B z \ \& \ x \neq_B z.$

CNs as Setoids: Copredication

- With these definitions, the desired semantics of our copredication cases are derived
- $Three(Book, \text{PHY} \bullet \text{INFO}, pm(j))$
- $\exists x, y, z : Book.D[\text{PHY}](x, y, z) \ \& \ D[\text{INFO}](x, y, z) \ \& \ pm(j, x) \ \& \ pm(j, y) \ \& \ pm(j, z)$
 - ▶ Note that this is achieved through defining the equivalence relation for dot-types by means of disjunction of both identity criteria and, then, we obtain double distinctness by negating the disjunction.

CNs as Setoids: Copredication

- Verbs plus adjectives in quantified copredication

- ▶ Consider the following example:

(23) John mastered three heavy books.

- ▶ The interpretation needed here: John mastered three informational objects that are also heavy as physical objects

- ★ Both the verb and the adjective have a word on the IC

- First step: adjectival modification

- ▶ $HBook = \Sigma(Book, heavy)$ or $\Sigma x:Book.heavy(x)$

CNs as Setoids: Copredication

- The interpretation we get:

$Three(HBook, PHY \bullet INFO, master(j))$

- Expanding:

$\exists x, y, z : HBook.D[PHY](x, y, z) \ \& \ D[INFO](x, y, z)$
 $\ \& \ master(j, x) \ \& \ master(j, y) \ \& \ master(j, z)$

Conclusions

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Conclusions

- MTTs as foundational languages for formal semantics
 - ▶ Formally, well-studied
 - ▶ Expressively adequate
 - ▶ Proof-theoretically specified, supporting effecting reasoning
- State of maturity of both MTT semantics and proof assistant technology
 - ▶ Use proof assistant technology and MTTs for formal verification and inference

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